CDC patterns in Apache Iceberg

Ryan Blue Trino Fest – June 2023



Scan for an Iceberg cheat sheet for Spark or Trino



Quick refresher

What's Iceberg?

Iceberg is an **open standard** for tables with **SQL behavior**

What's CDC?

Change Data Capture: As relational tables are modified, emit an update stream to keep copies in sync—capture changes to tables as they happen





Bank example

Bank accounts

- Account ID and balance
- Updated by primary key
- Layout and order configured

Goal: Keep accounts up-to-date using incoming transaction data

```
-- example table
CREATE TABLE accounts (
        account_id bigint,
        balance decimal(12, 2))
PARTITIONED BY (
        bucket(4, account_id))
```

-- set primary key fields
ALTER TABLE accounts
SET IDENTIFIER FIELDS account_id

-- configure write order/distribution ALTER TABLE accounts WRITE DISTRIBUTED BY PARTITION LOCALLY ORDERED BY account_id



Transaction data

Double entry bookkeeping

- Each transfer updates 2 accounts
- Total deposits should not change (transactional consistency)

Transaction source is flexible

- Kafka or kinesis stream
- Upstream table

transaction_id	account_id	amount
+ 	9	
1	8	435
2	2	-863
2	4	863
3	6	-530
+	+	+

SELECT sum(balance) AS total_deposits FROM accounts



Wants

- Direct writes single table
- Accurate historical record
- Time travel to any point
- Consistent within and across tables
- High volume, low latency
- Read-optimized
- Write-optimized
- Schema evolution



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Problems

- Lower latency \Rightarrow more work
- Write amplification
- Batch writes frequency
 - Double update problem
 - Transaction alignment/consistency
- Read requirements
 - Equality: delete *id=5*
 - Positional: delete *A.parquet, pos 11*



Storage trade-off

Direct writes

- One table, one write
- Increases write complexity
- Volume limit
- Double update problem

Change log table

- Historical record
- Time travel to any transaction
- Simple append-only writes
- High volume
- No direct reads, not optimized

Most important (and overlooked) decision



Change log pattern

Surprisingly effective with Trino!

- Track only changes
- Efficient writes, expensive reads
- Continuous time travel:
 WHERE transaction_id < ID

Tip: Handle UPSERT using SQL windows

```
-- store only account changes
CREATE TABLE account_updates (
    transaction_id bigint,
    account_id bigint,
    amount decimal(12, 2))
PARTITIONED BY (
    truncate(100000, transaction_id))
```

-- compute account value at query time
CREATE VIEW accounts AS
SELECT

account_id, sum(amount) AS balance FROM account_updates



MERGE pattern

- Direct write to an analytic table
- Uses position deletes (reads data!)
- Supports custom logic
 - Count duplicates
 - Consume any source data

```
-- squash multiple updates
WITH updates AS (
    SELECT
        account_id,
        sum(amount) AS amount
    FROM transactions
    GROUP BY account_id
)
```

```
MERGE INTO accounts a USING updates u
ON a.account_id = u.account_id
WHEN MATCHED THEN UPDATE
    SET a.balance = a.balance + u.amount
```



MERGE strategy trade-off

Lazy

merge-on-read

- Write only updates
- Low write amplification
- Defer work to read or compaction
- Example table: creates up to 8 files

Eager

copy-on-write

- Rewrite files as needed
- High write amplification
- Do work at write time for fast reads
- Example table: rewrites up to 4 files

Supported in Spark and Trino

Supported in Spark



Commit frequency trade-off

Faster

- Closer to real time
- Requires more maintenance
- Exacerbates the strategy trade-off!

Slower

- Higher latency for changes
- Reduces conflicts with services

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Flink UPSERT pattern

Update type trade-off

- Equality update
 - No reading needed
 - Cannot compact deltas
- Positional update
 - Requires locating rows
 - Can conflict with updates

Flink UPSERT is NOT recommended

- Inflexible
- Requires aggressive maintenance
- Doesn't sort data for efficiency
- Worst pattern in practice





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Reasons to use the change log pattern



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Reasons to use the MERGE pattern

- Use eager rewrites by default (copy-on-write)
- Use lazy rewrites for frequent updates



Hybrid pattern: MERGE + change log

Best of both patterns

- Land updates in change log table
 - Optimized for writes
 - Historical record, time travel
- Periodically MERGE
 - Simple reads
 - Separates concerns
- Optional view for read efficiency
 - Low data latency
 - Infrequent MERGE

Worst of both patterns

- Eager/lazy strategy trade-off
- Commit frequency trade-off
- Complex pipeline



Future work

Branches and tags

- Maintain change log in a branch
- Tag periodic MERGE results
- Use views to apply latest changes

New patterns

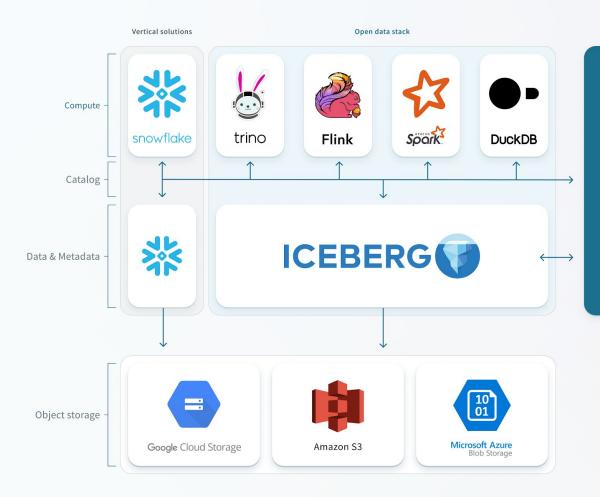
- LSM patterns
 - Equality updates with sorted data





Questions?

Thanks for attending! app.tabular.io/signup

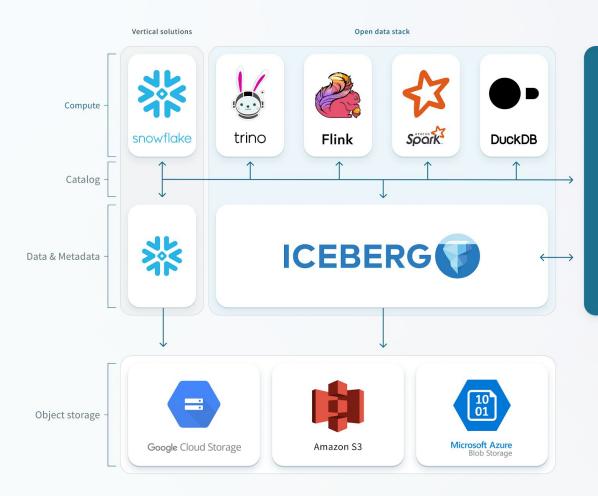








What is Tabular?









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