



Wvlet: A New Flow-Style Query Language

For Functional Data Modeling and Interactive Data Analysis

December 12, 2024
Trino Summit 2024

Taro L. Saito (Leo)
Senior Principal Engineer at Treasure Data

2024: 50th Anniversary of SQL

- SQL was born in 1974
 - SQL's early design choices paved the way for standardization and commercial adoption
- Since then, applications of SQL have expanded significantly
 - OLAP: From RDBMS to Trino, Hive, Spark, etc.
 - Embedded DB: SQLite (OLTP) -> DuckDB (OLAP with Parquet/Iceberg support)
 - SQL's user base has grown beyond DBA to include non-engineers

research and advances

DOI:10.1145/3648857

A discussion of the evolution of the database industry over the past half century, and why the relational database concepts introduced by E.F. Codd have proven to be so resilient over several decades.

BY DONALD CHAMBERLIN

50 Years of Queries

In the 1890 U.S. census. Punched cards had a successful 65-year product life until they were largely replaced by magnetic tapes in the 1950s. In the mid-20th century, data was typically stored on a magnetic tape and dedicated to a specific application. A tape might, for example, be used by an inventory-control application. Periodically, maybe once a week, the inventory-control job would read the tape sequentially, applying updates as it went along and producing a new, updated inventory tape. (As a college student in 1964, I had a summer job as a computer operator, running jobs like this.)

The advent of magnetic disks, introduced with the IBM RAMAC in 1956,¹ had a radical impact on how data was stored and processed. It was no longer necessary for applications to process data sequentially, since data items stored on disks could be accessed directly in any order. This gave rise to a new wave of innovation in how data should be organized on disk.

In the 1960s, a team of IBM engineers working on a NASA contract developed a disk-based information storage and retrieval system for use in the Apollo moon landing program.

2 key insights

- The relational data model, proposed by E.F. Codd in 1970, is the most widely used format for business data. Its practical feasibility was demonstrated in the 1970s by experimental prototypes at IBM Research and the University of California. The 1980s saw a proliferation of relational database products.
- SQL (later shortened to SQL) was designed in 1974 as a language for interacting with a relational database. It has been used widely by professional programmers. SQL was adopted as an ANSI Standard and by the availability of high-quality query engine implementations. Today, SQL remains the most widely used query language.

Current requirements for massive scale data processing are driving the evolution of SQL.

the 1890 U.S. census. Punched cards had a successful 65-year product life until they were largely replaced by magnetic tapes in the 1950s. In the mid-20th century, data was typically stored on a magnetic tape and dedicated to a specific application. A tape might, for example, be used by an inventory-control application. Periodically, maybe once a week, the inventory-control job would read the tape sequentially, applying updates as it went along and producing a new, updated inventory tape. (As a college student in 1964, I had a summer job as a computer operator, running jobs like this.)

The advent of magnetic disks, introduced with the IBM RAMAC in 1956,¹ had a radical impact on how data was stored and processed. It was no longer necessary for applications to process data sequentially, since data items stored on disks could be accessed directly in any order. This gave rise to a new wave of innovation in how data should be organized on disk.

In the 1960s, a team of IBM engineers working on a NASA contract developed a disk-based information storage and retrieval system for use in the Apollo moon landing program.

2 key insights

- The relational data model, proposed by E.F. Codd in 1970, is the most widely used format for business data. Its practical feasibility was demonstrated in the 1970s by experimental prototypes at IBM Research and the University of California. The 1980s saw a proliferation of relational database products.
- SQL (later shortened to SQL) was designed in 1974 as a language for interacting with a relational database. It has been used widely by professional programmers. SQL was adopted as an ANSI Standard and by the availability of high-quality query engine implementations. Today, SQL remains the most widely used query language.

Current requirements for massive scale data processing are driving the evolution of SQL.

A Relational Model of Data for Large Shared Data Banks

E. F. Codd
IBM Research Laboratory, San Jose, California

Future users of large data banks must be planning to know how the data is organized in the internal representation. A programmer must know the internal representation of the data.

IBM

TREASURE DATA

Trino at Treasure Data (2024)

- Treasure Data has been operating Trino (formerly Presto) as a service since 2014 (10th anniversary!)
 - 3+ million Trino SQL queries processed / day
 - 400+ trillion rows processed / day
 - 3+ billion S3 GET requests / day
 - Reduced from 10 billion requests / day (2023) by partition optimization
- In 2024, we completed the customer traffic migration from Presto (350) to Trino
 - See our migration/test methods in DBTest 2022 paper
- A lot of challenges in managing SQL and helping users (or LLM) write efficient SQL queries



trino



TREASURE
DATA

Journey of Migrating Millions of Queries on The Cloud

Taro L. Saito, Naoki Takezoe, Yukihiro Okada, Takako Shimamoto, Dongmin Yu, Suprith Chandrashekarachar, Kai Sasaki, Shohei Okumiya, Yan Wang, Takashi Kurihara, Ryu Kobayashi, Keisuke Suzuki, Zhenghong Yang, Makoto Onizuka*
Treasure Data
Mountain View, CA, USA

ABSTRACT

Treasure Data is processing millions of distributed SQL queries every day on the cloud. Upgrading the query engine service at this scale is challenging because we need to migrate all of the production queries of the customers to a new version while preserving the correctness and performance of the data processing pipelines. To ensure the quality of the query engines, we utilize our query logs to build customer-specific benchmarks and replay these queries with real customer data in a secure pre-production environment. To simulate millions of queries, we need effective minimization of test query sets and better reporting of the simulation results to proactively find incompatible changes and performance regression of the new version. This paper describes the overall design of our system and shares various challenges in maintaining the quality of the query engine service on the cloud.

ACM Reference Format:
Taro L. Saito, Naoki Takezoe, Yukihiro Okada, Takako Shimamoto, Dongmin Yu, Suprith Chandrashekarachar, Kai Sasaki, Shohei Okumiya, Yan Wang, Takashi Kurihara, Ryu Kobayashi, Keisuke Suzuki, Zhenghong Yang, Makoto Onizuka*
Treasure Data
Mountain View, CA, USA

As of February 2022, Treasure Data is processing more than 1.5 million SQL queries every day coming from 5,000+ users in various regions, including US, EU, Japan, Korea, etc. The number of processed records exceeds 100 trillion rows/day, which is equivalent to processing 1.2 billion rows/sec. Our job as a service provider is preserving the behavior of these SQL queries while maintaining the underlying query engine veritas up-to-date so that our customers can keep processing their data analysis pipelines without worrying about maintaining the data platform.

Maintaining A Data Platform on the Cloud. Treasure Data provides query engine services on the cloud by integrating open-source (OSS) distributed SQL query engines Trino (formerly known as Presto) [2] and Hive [1] on top of Amazon Web Service (AWS). Our target customers (e.g., marketers) are not dedicated system engineers, so they have a strong demand to have a data platform which is easy to use and has no need to maintain by themselves.

[cs.DB] 17 May 2022

DBTest '22, June 17, 2022, Philadelphia, PA, USA

Taro and Nasaki, et al.

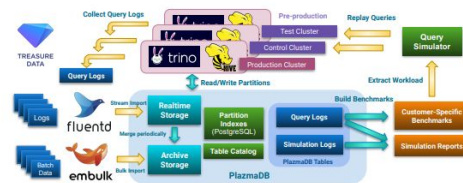


Figure 1: The architecture of Treasure Data, including PlazmaDB, query engines, and query simulator

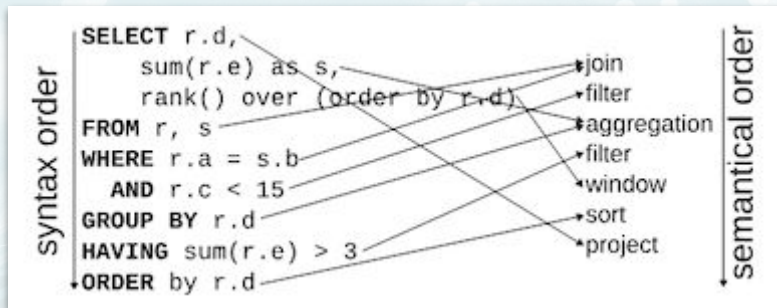
format of query logs changes over time. It is also important to extract typical workloads of individual customers to reduce the number of queries to test. Our approach for these problems is using our own service for collecting and analyzing query logs. An analysis of our historical query logs revealed that more than 97% of queries of our customers are recurrent ones, running hourly, daily, weekly, monthly, etc. By clustering such recurrent queries, we can significantly reduce the query set size fed to our query simulator. In addition, to protect the privacy of the customer data, the query simulator obfuscates the query results by embedding a checksum computation process and ensures the containment of all intermediate data within a pre-production environment, which is isolated from the production customer traffic.

Simulating customer-specific query patterns has enabled us to fill the gap of responsibilities between service providers like Treas-

ure Data and customers to test their Hyper in-memory DBMS [16]. Such static benchmarks, however, are not always a representative workload of customers. Recent approaches for testing customer-specific workloads include Snowtrail for testing Snowflake SQL queries [20] and DIAMetrics [5], which also covers non-SQL workloads in Google. A formal framework for constructing a representative workload from query logs is studied in [4], which requires the expert knowledge to select features to be tested. In this paper, we argue the importance of collecting query logs even for testing emerging OSS query engines using un supervised methods. In this regard, we shed lights to the system for collecting query logs and practical applications of these query logs for monitoring and testing query engines. Morphheus [13] reported a method for defining SLOs of query performance induced from large query graphs. In this paper, we present a simple approach for

What's Wrong with SQL?

- The syntactic order of SQL doesn't match the actual data flow
 - [A Critique of Modern SQL And A Proposal Towards A Simple and Expressive Query Language \(CIDR 24\)](#)
 - Even SQL experts find it challenging to debug nested queries



- Lacks essential software engineering features for managing many queries
 - No built-in support for reusing queries
 - No entry points for multi-query optimization
 - e.g., incremental processing and pipeline execution like dbt
 - No built-in debugging or testing capabilities

Wvlet: Modernizing 50-Year-Old SQL

- SQL has proven to be effective and useful
 - Its natural-language based syntax has gained widespread adoption
- Redesigned SQL to be more intuitive and functional
- Wvlet: a new flow-style query language for weaving data
 - Pronounced as weave-let
 - Queries start with “from” for intuitive data flow
- Functional
 - Wvlet queries are reusable and composable like functions
- Cross SQL engine support
 - Generates SQL for Trino, Hive, DuckDB, etc.
- Try Wvlet on your web browser <https://wvlet.org/>
 - No installation is required

```
ww> -- Scan a data file
from 'lineitem.parquet'
-- Apply a filtering condition
where l_quantity > 10.0
-- Aggregate by columns
group by l_returnflag, l_linestatus
-- Compute aggregated expressions
agg
  l_quantity.sum as qty,
  l_extendedprice.sum as price,
  l_discount.avg.round(4) as discount_avg,
-- Sort results
order by qty desc
-- Remove an unnecessary column
exclude discount_avg
-- Save query with an alias
select as stats
;
```

l_returnflag string	l_linestatus string	qty decimal(38,2)	price decimal(38,2)
N	O	732001.00	1026373433.61
R	F	365309.00	511907681.35
A	F	364109.00	509533926.17
N	F	8642.00	11948259.73

4 rows

ww> █

Command-line editor of Wvlet (ww)

Wvlet Playground wvlet.org

The screenshot displays the Wvlet Playground interface. On the left is a sidebar with a 'Playground' section and a list of 'Examples' including '00_sample.wv', 'add.wv', 'agg.wv', 'column_ops.wv', 'join.wv', 'list_schemas.wv', 'list_tables.wv', 'model.wv', 'model_with_param.wv', 'raw_sql.wv', 'scan.wv', 'shift_columns.wv', 'string_interpolation.wv', 'tpch_q1.wv', 'tpch_q2.wv', 'tpch_q3.wv', 'tpch_schema.wv', and 'values.wv'. The main area is split into three panels: 'Wvlet' (Query Syntax), 'Compiled SQL', and 'Preview'. The 'Wvlet' panel contains a SQL query with 14 lines of code. The 'Compiled SQL' panel shows the same query with additional metadata. The 'Preview' panel shows a table with 4 rows and 3 columns: 'l_linestatus', 'sum_qty', and 'sum_ext_price'. Below the table is a 'Query Result' button.

```
Wvlet 00_sample.wv
1 -- Scan a table
2 from lineitem
3 -- Add filtering condition
4 where l_quantity > 0.0
5 -- Grouping by keys
6 group by l_returnflag, l_linestatus
7 -- Add aggregation expressions
8 agg
9   l_quantity.sum as sum_qty,
10  l_extendedprice.sum as sum_ext_price
11 -- Remove unnecessary column
12 exclude l_returnflag
13 -- Sort
14 order by sum_qty desc
```

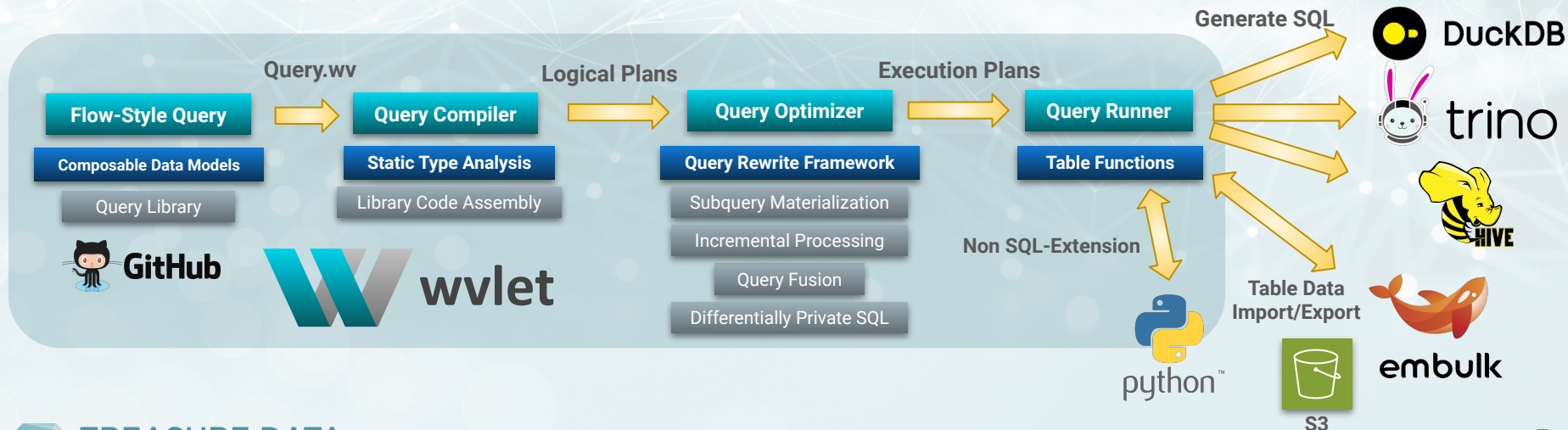
```
Compiled SQL
1 -- wvlet version=2024.9.13
  +33-7bdc7959, src=01JEPS6BY3DWGB0N50ZBFAX934.wv:14
2 select * from
3   (select l_linestatus, sum_qty, sum_ext_price from
4     (select l_returnflag as l_returnflag, l_linestatus as
5       l_linestatus, sum(l_quantity) as sum_qty, sum
6       (l_extendedprice) as sum_ext_price
7     from
8       (select * from lineitem
9         where l_quantity > 0.0)
9     group by l_returnflag, l_linestatus))
  order by sum_qty desc
```

l_linestatus string	sum_qty decimal(38,2)	sum_ext_price decimal(38,2)
0	76525100	107286230210
F	38144900	53459444535
F	38045600	53234821165
F	897100	1238480137

4 rows

Wvlet: Architecture

- From query (.wv) files, Wvlet compiler produces logical plans, execution plans, and SQL statements
- Logical Plans
 - Tree-representations of relational operators (e.g., scan, filter, projection)
 - Sort(Join(Filter(TableScan(...)), TableScan(...)))
- Execution Plans
 - A sequence of steps to execute SQL and other programs
 - ExecutePlan(ExecuteSQL(query), ExecuteTest(expr), ExecuteCommand(expr), ...)



Wvlet: Analyze As You Write

- Flow-style queries
 - Each line is a single operation
- Peek the data and schema at any point
 - A subquery becomes a range of lines in the query text
- Easy to extend
 - Queries can be reused for further analysis



Related Work: GoogleSQL Pipe Syntax (2024)

- [SQL Has Problems. We Can Fix Them: Pipe Syntax In SQL](#) (VLDB 2024)
- Extended SQL syntax with [pipe operator \(|>\)](#)
 - Available in ZetaSQL, an open-source SQL parser, used in Google products, e.g., BigQuery, F1, etc.
- Other flow-style languages:
 - PRQL, Microsoft Kusto, DryadLINQ, etc.

```
SELECT c_count, COUNT(*) AS custdist
FROM
  ( SELECT c_custkey, COUNT(o_orderkey) c_count
    FROM customer
    LEFT OUTER JOIN orders ON c_custkey = o_custkey
                        AND o_comment NOT LIKE '%unusual%packages%'
    GROUP BY c_custkey
  ) AS c_orders
GROUP BY c_count
ORDER BY custdist DESC, c_count DESC;
```

SQL



SQL Has Problems. We Can Fix Them: Pipe Syntax In SQL

Jeff Shute Google, Inc.	Shannon Bales Google, Inc.	Matthew Brown Google, Inc.	Jean-Daniel Browne Google, Inc.	Brandon Dolphin Google, Inc.
Romit Kudtarkar Google, Inc.	Andrey Litvinov Google, Inc.	Jingchi Ma Google, Inc.	John Morcos Google, Inc.	Michael Shen Google, Inc.
	David Wilhite Google, Inc.	Xi Wu Google, Inc.	Lulan Yu Google, Inc.	

sql-pipes-paper@google.com

ABSTRACT

SQL has been extremely successful as the de facto standard language for working with data. Virtually all mainstream database-like systems use SQL as their primary query language. But SQL is an old language with significant design problems, making it difficult to learn, difficult to use, and difficult to extend. Many have observed these challenges with SQL, and proposed solutions involving new languages. New language adoption is a significant obstacle for users, and none of the potential replacements have been successful enough to displace SQL.

In GoogleSQL, we've taken a different approach - solving SQL's problems by extending SQL. Inspired by a pattern that works well in other modern data languages, we added piped data flow syntax

SQL. Migrating away from existing SQL ecosystems is expensive and generally unappealing for users.

This paper presents a different approach. After describing the most critical problems with the SQL language, we present a solution - adding pipe-structured data flow syntax to SQL. This makes SQL more flexible, extensible and easy to use. This paradigm works well in other languages like Kusto's KQL [5] and in APIs like Apache Beam [1]. We show pipe syntax can be added to SQL, too, without removing anything, and while maintaining full backwards compatibility and interoperability.

In SQL, the standard clauses occur in one rigidly defined order. Expressing anything else requires subqueries or other workarounds. With pipe syntax, operations can be composed arbitrarily, in any

```
FROM customer
|> LEFT OUTER JOIN orders ON c_custkey = o_custkey
    AND o_comment NOT LIKE '%unusual%packages%'
|> AGGREGATE COUNT(o_orderkey) c_count
GROUP BY c_custkey
|> AGGREGATE COUNT(*) AS custdist
GROUP BY c_count
|> ORDER BY custdist DESC, c_count DESC;
```

SQL + Pipe Syntax

Wvlet: Relational Operators ([online reference](#))

- In Wvlet, all keywords must be lower-case letters
 - For consistency and reducing implementation efforts
- Same operators with SQL:
 - **from, select, join, where, order by, limit**, etc.
- Aggregation in Wvlet
 - **group by** k1, k2, ... **agg** ...
 - **agg** (aggregation expr), ..
 - Report group-by keys and aggregation expressions: k1, k2, ..., expr1, expr2, ...
 - **group by** k1, k2, ... **where** ...
 - equivalent to group by ... having ...

```
wv> from lineitem
      group by l_shipmode
      agg  _.count, l_quantity.sum;
```

l_shipmode string	count_star() long	sum(l_quantity) decimal(38,2)
FOB	8641	219565.00
SHIP	8482	217969.00
REG AIR	8616	219015.00
TRUCK	8710	223909.00
AIR	8491	216331.00
MAIL	8669	221528.00
RAIL	8566	217810.00

7 rows

Aggregation Query

```
wv> from ...      -- Scan the input data
where ...        -- Apply filtering conditions
where ...        -- [optional] Apply more filtering conditions
add ... as ...   -- Add new columns
transform ...    -- Transform a subset of columns
group by ...     -- Grouping rows by the given columns
agg ...          -- Add group aggregation expressions,
                 -- (e.g., _.count, _.sum)
where ...        -- Apply filtering conditions for groups
                 -- (e.g., HAVING clause in SQL)
exclude ...      -- Remove columns from the output
shift ...        -- Shift the column position to the left
select ...       -- Select columns to output
order by ...     -- Sort the rows by the given columns
limit ...        -- Limit the number of rows to output
```

Basic Flow of Wvlet Queries

Wvlet: Column At A Time Manipulation

- For minimizing your typing effort
- **add (expr) as ...**
 - Add a new column
- **rename (column) as (new name)**
- **exclude (column), ...**
- **shift (to left/right)? column,**
 - Reorder columns for readability
- Changing a single column or its order in SQL is not easy as we need to enumerate all columns

```
wv> from lineitem
  where l_returnflag = 'N'
  add s"${l_shipmode}:${l_shipinstruct}" as label
  shift label, l_returnflag
  limit 5;
```

label string	l_returnflag string	l_shipmode string	l_shipinstruct string	l_quantity
TRUCK:DELIVER IN PERSON	N	TRUCK	DELIVER IN PERSON	
MAIL:TAKE BACK RETURN	N	MAIL	TAKE BACK RETURN	
REG AIR:TAKE BACK RETURN	N	REG AIR	TAKE BACK RETURN	
AIR:NONE	N	AIR	NONE	
FOB:NONE	N	FOB	NONE	

5 rows

Wvlet

```
-- wvlet version:2024.9.9, src:01JBYY824NEXQ9BTC9RKYEREFQ.wv
select * from
  (select label, l_returnflag, l_shipmode, l_shipinstruct, l_quantity, l_tax, l_shipdate, l_suppkey, l_orderkey, l_comment,
  l_linestatus, l_receiptdate, l_linenum, l_extendedprice, l_partkey, l_discount, l_commitdate from
    (select *, concat(concat(concat(concat('', l_shipmode), ':'), l_shipinstruct), '') as label from
      (select * from lineitem
        where l_returnflag = 'N'))))
  limit 5 - (QueryExecutor.scala:319)
```

SQL

Wvlet: Additional Relational Operators

- Refined from SQL
 - **concat** = UNION ALL in SQL
 - Most of SQL users **wrongly use UNION (duplicate elimination)** where UNION ALL (concatenate two relations) is appropriate.
 - **dedup** = select distinct * (= duplicate elimination)
- Utility operators
 - **transform**: Update only subset of columns
 - **sample n**
 - **pivot**
 - Transform column values into individual columns
 - [trino#1206](#) (Remains open since 2019)
 - Wvlet can run multiple SQL queries, so it's relatively easy to implement static/dynamic pivot

```
wv> from lineitem
  where L_returnflag = 'N'
  limit 3
concat {
  from lineitem
  where L_returnflag = 'A'
  limit 3
}
shift L_returnflag
```

Wvlet

```
-- wvlet version:2024.9.9, src:01JBYX6C240NPZZ9ZR3Y0T9YA9.wv
select L_returnflag, L_shipmode, L_shipinstruct, L_quantity, L_tax, L
shipdate, L_suppkey, L_orderkey, L_comment, L_linestatus, L_receiptdat
e, L_linenum, L_extendedprice, L_partkey, L_discount, L_commitdate
from
(
  (select * from
    (select * from lineitem
     where L_returnflag = 'N')
   limit 3)
 union all
  (select * from
    (select * from lineitem
     where L_returnflag = 'A')
   limit 3))) - (QueryExecutor.scala:319)
```

SQL

Wvlet: Update Statements

- **save as**
 - A shorthand notation for:
 - DROP TABLE IF EXISTS tbl; CREATE TABLE tbl AS SELECT ...
 - Most frequently used pattern in Treasure Data
 - CREATE OR REPLACE TABLE AS SELECT
 - Available in DuckDB, [Trino 431](#) (Since Oct 2023)
 - For DuckDB backend, you can save query results as Parquet/JSON files
- **append to**
 - Almost same with INSERT INTO, but it clarifies append-only semantics (no overwrites)
 - Useful when using versioned tables like Iceberg/Delta Lake
- **delete**
 - Delete selected rows from the table

```
wv> from lineitem
  where l_returnflag = 'N'
  limit 10
  save as tmp_table;
```

Wvlet

```
Executing SQL:
create or replace table tmp_table as
-- wvlet version:2024.9.9, src:01JBYYKEV8CK8HTZEBZGAMYVZB.wv
select * from
(select * from lineitem
 where l_returnflag = 'N')
limit 10 - (QueryExecutor.scala:173)
```

SQL

Function Chaining via Dot Operator

- In SQL, every function is global
 - SUM(x), AVG(x), COUNT(*), ROUND(x, 1), ...
 - A lot of cursor movements are necessary to apply functions in SQL
- Wvlet supports chaining functions with dot operator as in modern programming languages:
 - e.g., x.sum, x.avg.round(2), _.count
 - _ (underscore) refers to the output from the previous operator

```
select
  l_shipmode,
  round(avg(l_quantity), 1),
  sum(l_quantity),
  count(*)
from lineitem
group by l_shipmode
;
```

SQL

```
wv> from lineitem
  group by l_shipmode
  agg
    l_quantity.avg.round(1),
    l_quantity.sum,
    _.count,
  ;
```

l_shipmode string	round(avg(l_quantity), 1) double	sum(l_quantity) decimal(38,2)	count_star() long
REG AIR	25.4	219015.00	8616
FOB	25.4	219565.00	8641
SHIP	25.7	217969.00	8482
MAIL	25.6	221528.00	8669
RAIL	25.4	217810.00	8566
AIR	25.5	216331.00	8491
TRUCK	25.7	223909.00	8710

7 rows

Wvlet

Wvlet Query = A Function for Building A Data Model

- In Wvlet, a query defined as a data model that can be reused at ease (Query Templates)
- Composable
 - You can add more relational operators (e.g., where, join) to data models
- Reusable
 - Models accept user parameters (= function arguments)

```
wv> -- define a reusable query for analyzing cost
| model service_cost =
|   from cost_components
|   where time.within('-3M/now')
|   where aws_account_owner != 'security'
|   add service_name = coalesce(user_service, cost_category)
|   add td_time_string(time, 'M!') as month
|   pivot on month in ('2024-06', '2024-07', '2024-08')
|   group by service_name
|   agg cost.sum.round(1)
| end;
```

Reuse Queries



```
wv> from service_cost
| where service_name like 'trino%';
```

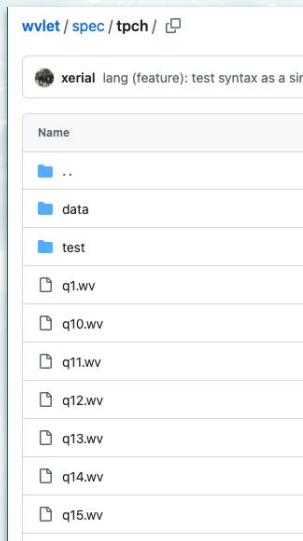
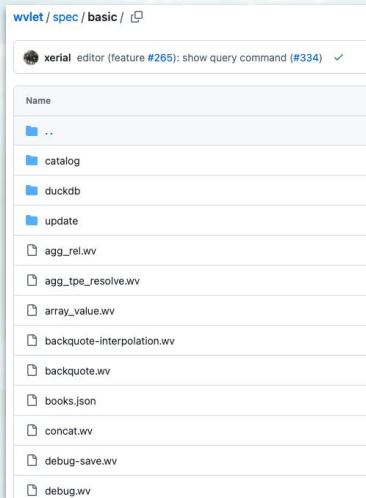
service_name string	2024-06 double	2024-07 double	2024-08 double
trino-coordinator	3415.3	3983.4	4402.7
trino-worker	33776.7	58363.2	78516.8
2 rows			

```
wv> from service_cost
| where service_name like 'hive%';
```

service_name string	2024-06 double	2024-07 double	2024-08 double
hivep	174.9	164.4	188.4
hive-query-simulator	39.6	26.9	28.6
2 rows			

Testing Queries

- Wvlet Test Syntax
 - Useful for verifying resulting schema and results
- Wvlet is tested with Wvlet
 - Wvlet spec queries
 - Covers all 22 TPC-H queries



```
person.json
[
  {"id":1, "name": "alice", "age": 10 },
  {"id":2, "name": "bob", "age": 24 },
  {"id":3, "name": "clark", "age": 40 }
]
```

```
from 'person.json'
-- Test the query size and schema
test _size should be 3
test _columns should be ['id', 'name', 'age']
test _columns should contain 'name'
test _columns should not contain 'address'

-- Test the query output. Surrounding white spaces will be trimmed:
test _output should be ""

  id | name | age
  long | string | long
  ---|---|---
  1 | alice | 10
  2 | bob   | 24
  3 | clark | 40

  3 rows

""""
```


Debugging Queries

- Debug operator can be used for checking intermediate query results
- ExecutionPlanner generates multiple execution paths for debug and regular query evaluation
 - Debug path: Query before debug statement -> debug query
 - Regular path: Query without debug statement

```
[  
  {"id":1, "name": "alice", "age": 10 },  
  {"id":2, "name": "bob", "age": 24 },  
  {"id":3, "name": "clark", "age": 40 }  
]
```

```
ww> from 'person.json'  
  debug {  
    where id = 1  
    save as debug_result_0123  
  }  
  where id = 2  
  
  from debug_result_0123  
  test ..size should be 1  
  test ..rows should be [[1, "alice", 10]]
```

debug (line:9): test ..rows should be [[1, "alice", 10]]

id long	name string	age long
2	bob	24
1 rows		

Regular Query Result

id long	name string	age long
1	alice	10
1 rows		

Debug Query Result

[ok]: 1 was equal to 1 (01JAXGC09FFAZ0ZNHAWWZSBJ4.wv:8:13)

Test Result

[ok]: [[1, alice, 10]] was equal to [[1, alice, 10]] (01JAXGC09FFAZ0ZNHAWWZSBJ4.wv:9:13)

Wvlet Internals: Compiler

- Compiler phases: Parser -> SymbolLabeler -> TypeResolver -> Rewriter -> ExecutionPlanner -> GenSQL
 - CompilationUnit holds the source text, untyped logical plan, typed logical plan, execution plan, etc.
- With [AirSpec](#) testing library, you can adjust the log levels of individual components
 - -L (class name pattern)=(log level) option

```
wv> from 'books.json'  
  group by author  
  agg _.count;
```

Query (.wv)



```
sbt:wvlet> --runner/testOnly *BasicSpec -- -l debug -L *Scanner=debug -L *GenSQL=trace agg_rel  
[info] compiling 1 Scala source to /Users/leo/work/git/wvlet/wvlet-api/.jvm/target/scala-3.3.4/classes ...  
BasicSpec:  
2024-10-23 13:29:36.786-0700 debug [WvletScanner] Indented(0,null) [ 0: 4] FROM: from - (WvletScanner.scala:221)  
2024-10-23 13:29:36.803-0700 debug [WvletScanner] Indented(0,null) [ 5:12] STRING_LITERAL: books.json - (WvletScanner.scala:221)  
2024-10-23 13:29:36.804-0700 debug [WvletScanner] Indented(0,null) [18: 5] GROUP: group - (WvletScanner.scala:221)  
2024-10-23 13:29:36.804-0700 debug [WvletScanner] Indented(0,null) [24: 2] BY: by - (WvletScanner.scala:221)  
2024-10-23 13:29:36.804-0700 debug [WvletScanner] Indented(0,null) [27: 6] IDENTIFIER: author - (WvletScanner.scala:221)  
2024-10-23 13:29:36.805-0700 debug [WvletScanner] Indented(0,null) [34: 3] AGG: agg - (WvletScanner.scala:221)  
2024-10-23 13:29:36.805-0700 debug [WvletScanner] Indented(0,null) [38: 1] UNDERSCORE: - - (WvletScanner.scala:221)  
2024-10-23 13:29:36.805-0700 debug [WvletScanner] Indented(0,null) [39: 1] DOT: . - (WvletScanner.scala:221)  
2024-10-23 13:29:38.367-0700 trace [GenSQL] IDENTIFIER: count - (WvletScanner.scala:221)  
[plan] runner / Test / testOnly is  
[Agg [34..45]] <t167(author:string, _:array<t123(id:string, title:string, author:string, year:long, genre:string)>>) => <t168(author:string, <empty>:??)?>  
  - author:string := author  
  - <empty>:?? := sql"count(*)"  
  [GroupBy [18..33]] <t123(id:string, title:string, author:string, year:long, genre:string)> => <t167(author:string, _:array<t123(id:string, title:string, author:string, year:long, genre:string)>>)  
    - author  
    [JSONFileScan [5..17]] => <t123(id:string, title:string, author:string, year:long, genre:string)>  
[SQL]  
-- wvlet version:2024.9.8+i-c700bcaa, src:agg_rel.wv  
select author as author, count(*)  
from 'spec/basic/books.json'  
group by author - (GenSQL.scala:87)  
2024-10-23 13:29:38.422-0700 debug [BasicSpec]
```

Tokens

LogicalPlan
(After typing)

Generated SQL

Query Result

author string	count_star long
J.D. Salinger	2
F. Scott Fitzgerald	2
Harper Lee	2
John Steinbeck	1

4 rows
- spec:basic:agg_rel.wv 1.79s

- (SpecRunner.scala:52)

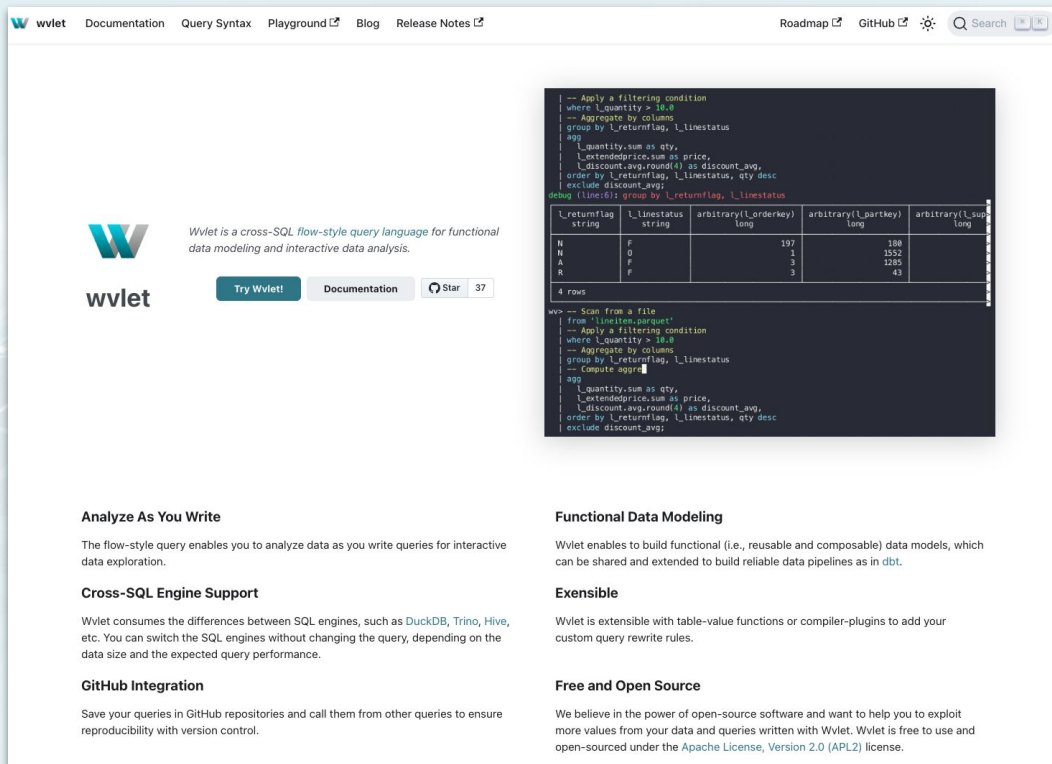


Wvlet Development Roadmap

- Roadmap is maintained at [GitHub Project](#) page.
 - Planning milestones for about every 3 months
 - Versions will be YYYY.(milestone month).(patch) (See [#170](#) for the versioning scheme)
- Release 2024.9
 - Flow-style query language design and compiler
 - Including scanner, parser, typer, tree rewrite framework, execution planner, SQL generator, etc.
 - DBMS Connector (DuckDB, Trino)
 - ww: Interactive command-line editor (REPL)
 - Installer (Homebrew)
 - Web UI, Playground (Monaco Editor with DuckDB-Wasm)
- Release 2025.1
 - Model management
 - Generate dependent model materialization plan, like dbt
 - GitHub integration
 - Compiler plugins for advance optimization
 - Incremental processing, query fusion, etc.
 - Language SDKs
 - Python, Rust, C/C++, Java, etc.
 - SQL to Wvlet converter

Summary

- Designed Wvlet, a new open-source flow-style query language
- Addresses challenges in 50-year-old design of SQL
 - Leverage good sides of SQL
 - natural-language like syntax
 - Intuitive syntax for data flow
 - Reusability
 - Extensibility
- Website: <https://wvlet.org/>
- GitHub: <https://github.com/wvlet/wvlet>



The screenshot shows the Wvlet website with a navigation bar containing links for Documentation, Query Syntax, Playground, Blog, and Release Notes. A search bar is located in the top right corner. The main content area features the Wvlet logo and a description: "Wvlet is a cross-SQL flow-style query language for functional data modeling and interactive data analysis." Below this is a "Try Wvlet!" button, a "Documentation" link, and a "Star 37" button. The page is divided into three columns of text:

- Analyze As You Write**: The flow-style query enables you to analyze data as you write queries for interactive data exploration.
- Cross-SQL Engine Support**: Wvlet consumes the differences between SQL engines, such as DuckDB, Trino, Hive, etc. You can switch the SQL engines without changing the query, depending on the data size and the expected query performance.
- GitHub Integration**: Save your queries in GitHub repositories and call them from other queries to ensure reproducibility with version control.
- Functional Data Modeling**: Wvlet enables to build functional (i.e., reusable and composable) data models, which can be shared and extended to build reliable data pipelines as in dbt.
- Extensible**: Wvlet is extensible with table-value functions or compiler-plugins to add your custom query rewrite rules.
- Free and Open Source**: We believe in the power of open-source software and want to help you to exploit more values from your data and queries written with Wvlet. Wvlet is free to use and open-sourced under the Apache License, Version 2.0 (APL2) license.

The right side of the screenshot displays a code editor with a Wvlet query and its output. The query is as follows:

```
-- Apply a filtering condition
where L_quantity > 10.0
-- Aggregate by columns
group by L_returnflag, L_linestatus
agg
  L_quantity, sum as qty,
  L_extendedprice, sum as price,
  L_discount, avg, round(4) as discount_avg,
order by L_returnflag, L_linestatus, qty desc
exclude discount_avg;
debug (line(6)): group by L_returnflag, L_linestatus
```

The output is a table with 4 rows:

L_returnflag	L_linestatus	arbitrary(L_orderkey)	arbitrary(L_partkey)	arbitrary(L_sup
string	string	long	long	long
N	F	197	180	
N	D	3	352	
A	F	3	1285	
R	F	3	43	

Below the table, another code block shows a query for scanning a file:

```
wv -- Scan from a file
| from "L_lines.parquet"
| -- Apply a filtering condition
| where L_quantity > 10.0
| -- Aggregate by columns
| group by L_returnflag, L_linestatus
| -- Compute aggregate
| agg
  L_quantity, sum as qty,
  L_extendedprice, sum as price,
  L_discount, avg, round(4) as discount_avg,
| order by L_returnflag, L_linestatus, qty desc
| exclude discount_avg;
```



Appendix

Wvlet Internals: Standard Library

- Wvlet defines standard data types and functions, which define how to generate SQL
 - example: `x.to_int.round(1) => cast(x as bigint).round(1) => round(cast(x as bigint),1)`

```
package wvlet.standard

type int:
  def to_int: int = sql"cast(${this} as bigint)"
  def to_long: long = sql"cast(${this} as bigint)"
  def to_float: float = sql"cast(${this} as double)"
  def to_double: double = sql"cast(${this} as double)"
  def to_boolean: boolean = sql"cast(${this} as boolean)"
  def to_string: string = sql"cast(${this} as varchar)"

  def or_else(other:int): int = sql"coalesce(${this},${other})"

  def round(decimal:int=0): double = sql"round(${this},${decimal})"

  def in(v:any*): boolean = sql"${this} in (${v})"
  def not_in(v:any*): boolean = sql"${this} not in (${v})"

  def between(l:int, r:int): boolean = sql"${this} between ${l} and ${r}"
end
```

Wvlet Internals: DBMS Specific Functions

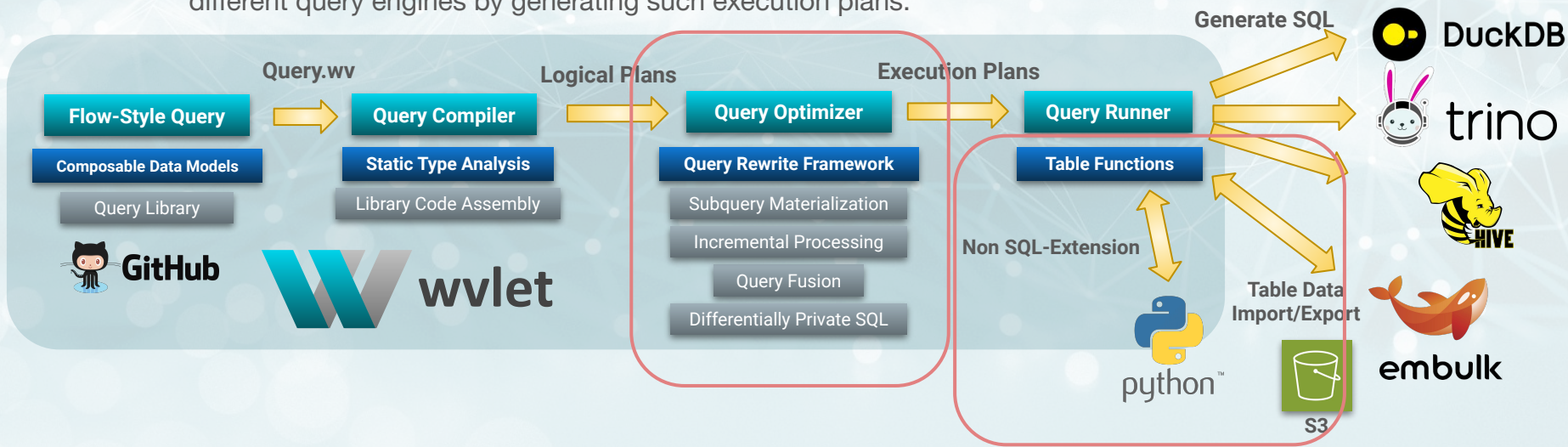
- Function can have different implementations depending on the target databases (Trino/Hive/DuckDB, etc.)
 - x.count_approx_distinct
 - Trino: approx_distinct(x)
 - DuckDB: approx_count_distinct(x)
- Works for consuming the differences between SQL dialects and UDFs

```
-- An array created by 'group by' clause
type array[A]:
  def length: int = sql"length(${this})"
  def size: int = sql"length(${this})"
  def get(index: int): A = sql"${this}[${index}]"
  def count: int = sql"count(*)"
  def count_distinct: int = sql"count(distinct ${this})"
  def count_if(cond:boolean): int = sql"count_if(${cond})"

-- Fast and memory-efficient approximate counting of distinct elements
def count_approx_distinct in trino: int = sql"approx_distinct(${this})"
def count_approx_distinct in duckdb: int = sql"approx_count_distinct(${this})"
```


Extending Wvlet

- **Compiler plugins** (to be designed [#185](#))
 - Add a custom rule set to optimize logical plans and execution plans.
 - Optimization rules: [Query fusion](#) optimization, [subquery materialization](#), [incremental processing](#), etc.
 - Security rules (e.g., forbidding local file access operators at the cloud environment)
- **Table functions** receive table-value data and output table value data
 - Not limited to SQL, we will be able to invoke ML algorithms, calling Embulk, or issuing SQL queries to different query engines by generating such execution plans.



Writing A Compiler From Scratch? Leverage LLMs

- Developing a compiler is a challenge with known difficulties and known solutions
 - = solved problem
- By leveraging LLM-based tools like GitHub Copilot, we can accelerate the development
- LLMs excel at utilizing well-known solutions from open-source code.

```
def parseJoin: Join =
  val t = scanner.lookAhead()
  t.token match
  case WvletToken.CROSS =>
    consume(WvletToken.CROSS)
    consume(WvletToken.JOIN)
    val right = relationPrimary()
    Join(JoinType.CrossJoin, EmptyRelation(t.span), right, NoJoinCriteria, spanFrom(t))
  case WvletToken.JOIN =>
    consume(WvletToken.JOIN)
    val right = relationPrimary()
    val joinOn = joinCriteria()
    Join(JoinType.InnerJoin, EmptyRelation(t.span), right, joinOn, spanFrom(t))
  case WvletToken.LEFT | WvletToken.RIGHT | WvletToken.INNER | WvletToken.FULL =>
    val joinType =
      t.token match
      case WvletToken.LEFT =>
        JoinType.LeftOuterJoin
      case WvletToken.RIGHT =>
        JoinType.RightOuterJoin
      case WvletToken.INNER =>
        JoinType.InnerJoin
      case WvletToken.FULL =>
        JoinType.FullOuterJoin
      case _ =>
        unexpected(t)
    consume(t.token)
    consume(WvletToken.JOIN)
    val right = relationPrimary()
    val joinOn = joinCriteria()
    Join(joinType, EmptyRelation(t.span), right, joinOn, spanFrom(t))
  case _ =>
    unexpected(t)
```

Indicate what you want to code

GitHub Copilot suggests code in a second

Wvlet is easy to learn, even for Large Language Models (LLMs)

- An LLM Agent, created from Wvlet documentation, can successfully convert TPC-H SQL queries into Wvlet syntax
 - Wvlet: Query Syntax documentation ([link](#))

Query Syntax

References

- Quick Start
- Relational Operators
- Expressions
- Metadata Functions
- Data Models

Introduction

Wvlet is a query language designed to be more human-readable and easier to write than SQL. If you already familiar to SQL, you will find it's easy to learn the syntax of wvlet as there are a lot of similarities between wvlet and SQL. Even if you are new to SQL, no worries! You can start learning wvlet from scratch. If you know about DataFrame in Python, it will help you understand the wvlet query language as chaining relational operators in the flow-style is quite similar to using DataFrame API.

Wvlet queries start with `from` keyword, and you can chain multiple relational operators to process the input data and generate the output data. The following is a typical flow of chaining operators in a wvlet query:

```
from ... -- Scan the input data
where ... -- Apply filtering conditions
where ... -- Optionally, apply more filtering conditions
agg ... as ... -- Add new columns
transform ... -- Transform a subset of columns
group by ... -- Grouping rows by the given columns
agg ... -- Add more aggregation expressions, e.g., count, sum
where ... -- Apply filtering conditions for groups (e.g., having clause in SQL)
exclude ... -- Remove columns from the output
shift ... -- Shift the column position to the left
select ... -- Select columns to output
order by ... -- Sort the rows by the given columns
limit ... -- Limit the number of rows to output
```

Unlike SQL, whose queries always must follow the `SELECT ... FROM ... WHERE ... GROUP BY ... ORDER BY ... LIMIT ...` structure, wvlet uses the **flow-style syntax** to match the syntax order with the data processing order as much as possible to facilitate more intuitive query writing.

Some operators like `agg`, `transform`, `agg`, `exclude`, `shift`, etc. are not available in the standard SQL, but these new operators have been added for reducing the amount of code and making the query more readable and easier to compose. Eventually, these operators will be translated into the equivalent SQL syntax.



You

Write some TPC-H queries in Wvlet



Agent

Certainly! I'll provide you with some examples of TPC-H queries translated into ad-hoc queries. Here are a few examples:

1. Query 1 - Pricing Summary Report:

```
from lineitem
where l_shipdate <= date('1998-12-01') - interval '90' day
group by
  l_returnflag,
  l_linestatus
agg
  sum_qty = _.l_quantity.sum,
  sum_base_price = _.l_extendedprice.sum,
  sum_disc_price = (_.l_extendedprice * (1 - _.l_discount)).sum,
  sum_charge = (_.l_extendedprice * (1 - _.l_discount) * (1 + _.l_tax)).sum,
  avg_qty = _.l_quantity.ave,
  avg_price = _.l_extendedprice.ave,
  avg_disc = _.l_discount.ave,
  count_order = _.count
order by
  l_returnflag,
  l_linestatus
```

```
SELECT
  l_returnflag,
  l_linestatus,
  sum(l_quantity) AS sum_qty,
  sum(l_extendedprice) AS sum_base_price,
  sum(l_extendedprice * (1 - l_discount)) AS sum_disc_price,
  sum(l_extendedprice * (1 - l_discount) * (1 + l_tax)) AS sum_charge,
  avg(l_quantity) AS avg_qty,
  avg(l_extendedprice) AS avg_price,
  avg(l_discount) AS avg_disc,
  count(*) AS count_order
FROM
  lineitem
WHERE
  l_shipdate <= CAST('1998-09-02' AS date)
GROUP BY
  l_returnflag,
  l_linestatus
ORDER BY
  l_returnflag,
  l_linestatus;
```

Original TPC-H
Q1 SQL



wv: Interactive Query Editor

- Available with [brew install wvlet/wvlet/wvlet](#)
 - Supports Trino, DuckDB via profile settings
- Shortcut keys for checking intermediate schema and results
 - ctrl-j, ctrl-d (describe the schema at the line)
 - ctrl-j, ctrl-t (test run the subquery upto the line)
 - ctrl-j, ctrl-r (run the whole query)

```
wv> from cost_components
  | where time.within('-3M/now')
  | where aws_account_owner != 'security'
  | add service_name = coalesce(user_service, cost_category)
  | -- type ctrl+j ctrl+t (test run) to debug the intermediate query result
  | add td_time_string(time, 'M!') as month
  | pivot on month in ('2024-06', '2024-07', '2024-08')
  | group by service_name
  | agg cost.sum.round(1)
  | order by `2024-08` desc
  | limit 10;
```

debug (line:6): add td_time_string(time, 'M!') as month

cost_category string	aws_account_code string	aws_account_owner string	user_service string
Production	td	engineering	waf-event-collector3
Production	kitchen	engineering-lead	core-eks-app-ronald-testing
Non Production	td	engineering	waf-event-collector3
Production	kitchen	engineering-lead	core-eks-app-scp-foo
Non Production	td	engineering	core-eks-app-scp-foo
Production	kitchen	engineering-lead	waf-event-collector3
Non Production	td	engineering	core-eks-app-scp-foo
Production	kitchen	engineering-lead	waf-event-collector3
Non Production	td	engineering	core-eks-app-scp-foo
Production	td	engineering	waf-event-collector3

Subquery Result

```
describe (line:4): where aws_account_owner != 'security'
```

column_name string	column_type string
cost_category	string
aws_account_code	string
aws_account_owner	string
user_service	string
service_owner_team	string
user_site	string
user_cluster	string
customer_id	string
plan_type	string
resource_product_region	string
resource_product_code	string
resource_usage	string
net_effective_cost	double
effective_cost	double
public_cost	double
cost	double
user_stage	string
time	long

18 rows

```
wv> from cost_components
  | where time.within('-3M/now')
  | -- type ctrl+j ctrl+d (describe) here to show the intermediate schema
  | where aws_account_owner != 'security'
  | add service_name = coalesce(user_service, cost_category)
  | add td_time_string(time, 'M!') as month
  | pivot on month in ('2024-06', '2024-07', '2024-08')
  | group by service_name
  | agg cost.sum.round(1)
  | order by `2024-08` desc
  | limit 10;
```

Intermediate Table Schema

Design Philosophy of Wvlet

- Link: <https://wvlet.org/wvlet/docs/internal/design>
- Matching syntax order with data flow
 - Typing from left-to-right, top-to-bottom order
- Use only lowercase keywords
 - Forbid mixing SELECT and select
- One operation at a time
 - SQL's SELECT operator is too powerful
 - Breaking down SELECT statement into simple operators (add, transform, exclude, shift, etc)
- Human and machine friendly-syntax
 - Query lines form subqueries
 - Trailing comma support is essential to correctly extract subqueries in an executable form
 - Enable in-query data/schema inspection

```
SELECT
  sum(c1),
  -- Rename c2 with an alias
  c2 as c2_new,
  -- skip c3 for exclusion
  -- Add a new computed column
  c4 + c5 as c101,
  -- Shift c6 and c7 to the end
  c8,
  ...
  ...,
  c100,
  c6,
  c7,
FROM tbl
```

SQL

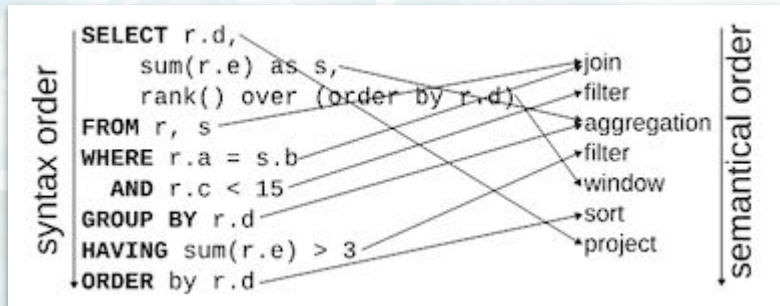


```
wv> from tbl
  -- Add a simple aggregation
  add c1.sum
  -- Rename c2 with an alias
  transform c2 as c2_new
  -- Remove c3 from the result
  exclude c3
  -- Add a new computed column
  add c4 + c5 as c101
  -- Shift c6 and c7 to the end
  shift to right c6, c7
```

Wvlet

Related Work: SaneQL (Thomas Neumann, et al. 2024)

- The syntactic order of SQL doesn't match with the semantic order of data processing
- Redesigned SQL by using a function chain syntax
 - [A Critique of Modern SQL And A Proposal Towards A Simple and Expressive Query Language \(CIDR 24\)](#)



```
let min_cost_for_part(p_partkey) :=  
    partsupp.filter(ps_partkey=p_partkey)  
    .join(supplier, s_suppkey=ps_suppkey)  
    .join(nation, s_nationkey=n_nationkey)  
    .join(region.filter(r_name='EUROPE'),  
          n_regionkey=r_regionkey)  
    .aggregate(min(ps_supplycost)),  
part  
    .filter(p_size = 15 && p_type.like('%BRASS'))  
    .join(partsupp, p_partkey = ps_partkey)  
    .join(supplier, s_suppkey = ps_suppkey)  
    .join(nation, s_nationkey = n_nationkey)  
    .join(region.filter(r_name='EUROPE'),  
          n_regionkey=r_regionkey)  
    .filter(ps_supplycost=min_cost_for_part(p_partkey)) ...
```

Related Work: PRQL (Prequel)

- Pipelined Relational Query Language (Since 2022) <https://prql-lang.org/>
- Written in Rust, compiling PRQL to SQL
 - Syntax is a bit far from SQL, more like DataFrame API or programming language syntax

PRQL

```
from tracks
filter artist == "Bob Marley" # Each line transforms the previous result
aggregate {                   # `aggregate` reduces each column to a value
  plays    = sum plays,
  longest  = max length,
  shortest = min length,      # Trailing commas are allowed
}
```

SQL

```
SELECT
  COALESCE(SUM(plays), 0) AS plays,
  MAX(length) AS longest,
  MIN(length) AS shortest
FROM
  tracks
WHERE
  artist = 'Bob Marley'
```

